

# COMPUTER SCIENCE & INFORMATION TECHNOLOGY

## Algorithms



Comprehensive Theory  
*with Solved Examples and Practice Questions*



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## **Algorithms**

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## **EDITIONS**

First Edition : 2015  
Second Edition : 2016  
Third Edition : 2017  
Fourth Edition : 2018  
Fifth Edition : 2019  
Sixth Edition : 2020  
Seventh Edition : 2021  
Eighth Edition : 2022

**Ninth Edition : 2023**

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# Algorithms

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## Goal of the Subject

This course provides an introduction to mathematical modeling of computational problems. It covers the common algorithms, algorithmic paradigms, and data structures used to solve these problems. The course emphasizes the relationship between algorithms and programming, and introduces basic performance measures and analysis techniques for these problems.

# Algorithms

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## INTRODUCTION

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In this book we tried to present the algorithms in a most simplified way. Each topic required for GATE is crisply covered with illustrative examples and each chapter is provided with Student Assignment at the end of each chapter so that the students get the thorough revision of the topics that he/she had studied. This subject is carefully divided into six chapters as described below:

**Chapter 1: Asymptotic Analysis of Algorithms:** In this chapter we discuss the asymptotic notations to represent the average, worst and best cases and we also learn how to identify time complexity of given algorithm.

**Chapter 2: Recurrence Relations:** In this chapter we study the three methods to study solve recurrence relations. The three methods are substitution method, master theorem and recursion tree method.

**Chapter 3: Divide Conquer:** In this chapter we discuss the various algorithms that can be implemented using divide and conquer paradigm namely merge and quick sort, Strassen's matrix multiplications. We also discuss the other sorting techniques.

**Chapter 4: Greedy Algorithms:** In this chapter we discuss the graph representations, algorithms' that require local optimizations at each step. Those algorithms include job scheduling, fractional Knapsack, Prim's and Kruskal's algorithms for MST's and other algorithms for shortest paths.

**Chapter 5: Graph based Algorithms:** In this chapter we discuss the BFS and DFS traversals. We also discuss the topological sorting algorithm.

**Chapter 6: Dynamic Programming:** In this chapter we discuss the algorithms for which greedy fails to give correct solution. We will discuss matrix chain multiplication, Floyd Warshall's algorithm, 0/1 Knapsack, longest common sequences and other algorithms.



# Asymptotic Analysis of Algorithms

## 1.1 Need for Performance Analysis

There are many important things that should be taken care of, like user friendliness, modularity, security, maintainability, etc. Why to worry about performance?

The answer to this is simple — we can have all the above things only if we have performance.

**Goal:** To write an algorithm which takes minimum time irrespective of the machine whether it is supercomputer or normal desktop.

### Choosing the Best Algorithm

One naive way of doing this is – implement both the algorithms and run the two programs on your computer for different inputs and see which one takes less time. There are many problems with this approach for analysis of algorithms.

1. It might be possible that for some inputs, first algorithm performs better than the second. And for some inputs second performs better.
2. It might also be possible that for some inputs, first algorithm perform better on one machine and the second works better on other machine for some other inputs.

ASYMPTOTIC ANALYSIS is the big idea that handles above issues in analyzing algorithms. In Asymptotic Analysis, we evaluate the performance of an algorithm in terms of input size (we don't measure the actual running time). We calculate, how does the time (or space) taken by an algorithm increases with the input size.

For example, let us consider the search problem (searching a given item) in a sorted array. One way to search is Linear Search (order of growth is linear) and other way is Binary Search (order of growth is logarithmic). To understand how Asymptotic Analysis solves the above mentioned problems in analyzing algorithms, let us say we run the Linear Search on a fast computer and Binary Search on a slow computer. For small values of input array size  $n$ , the fast computer may take less time. But, after certain value of input array size, the Binary Search will definitely start taking less time compared to the Linear Search even though the Binary Search is being run on a slow machine. The reason is the **order of growth** of Binary Search with respect to input size is logarithmic while the order of growth of Linear Search is linear. So the machine dependent constants can always be ignored after certain values of input size.

Asymptotic Analysis is not perfect, but that's the best way available for analyzing algorithms. For example, say there are two sorting algorithms that take  $1000 n \log n$  and  $2 n \log n$  time respectively on a machine. Both of these algorithms are asymptotically same (order of growth is  $n \log n$ ). So, With Asymptotic Analysis, we can't judge which one is better as we ignore constants in Asymptotic Analysis. Also, in asymptotic analysis, we always talk about input sizes larger than a constant value. It might be possible that those large inputs are never given to your software and an algorithm which is asymptotically slower always performs better for your particular situation. So, you may end up choosing an algorithm that is asymptotically slower but faster for your software.

#### Types of Asymptotic analysis:

1. Apriori analysis of algorithm.
2. Apostiari analysis of algorithms.

Apriori analysis	Apostiari analysis
1. It means we do analysis (space and time) of an algorithm prior to running it on specific system.	1. It means we do analysis of algorithm only after running it on system.
2. It does not depends on system.	2. It directly depends on system and changes from system to system
3. It provides approximate answer.	3. It gives exact answer.

## 1.2 Worst, Average and Best Cases

We can have three cases to analyze an algorithm:

1. Worst Case
2. Average Case
3. Best Case

Let us consider the following implementation of Linear Search.

```
#include <stdio.h>
// Linearly search x in arr []. If x is present then return the index,
// otherwise return -1
int search(int arr[], int n, int x)
{
    int i;
    for (i=0; i<n; i++)
    {
        if (arr[i] == x)
            return i;
    }
    return -1;
}
/* Driver program to test above functions*/
int main()
{
    int arr[] = {1, 10, 30, 15};
    int x = 30;
    int n = sizeof(arr)/sizeof(arr[0]);
    printf("%d is present at index %d", x, search(arr, n, x));
    getchar();
    return 0;
}
```



### 1.2.1 Worst Case Analysis

- In worst case analysis, we shall usually concentrate on finding only the **worst-case running time**, that is, the longest running time for any input of size  $n$ .
- The worst case analysis is an upper bound on the running time for any input. Knowing it gives us a guarantee that the algorithm will never take any longer.

**Example:** In the program given in section 1.2. That is for **linear search**, the worst case happens when the element to be searched is not present in the array (or) it is the last element, the search function compares it with all the elements of `arr[ ]` one by one. Therefore the worst case time complexity of linear search would be  $\theta(n)$ .

### 1.2.2 Average Case Analysis

- In average case analysis, we take all possible inputs and calculate computing time for all of the inputs.
- Considering all the cases, calculate the sum, and divide the sum by total number of inputs.

**Example:** For the linear search problem, let us assume that all cases are uniformly distributed (including the case of  $X$  not being present in array).

Considering all possible inputs:

i.e.,    when element found at 1<sup>st</sup> position —  $\theta(1)$   
           when element found at 2<sup>nd</sup> position —  $\theta(2)$   
           :  
           :  
           :  
           when element found at  $n^{\text{th}}$  position —  $\theta(n)$

$$\text{Sum} = \sum_{i=1}^n \theta(i)$$

$$\text{Average case time} = \frac{\sum_{i=1}^n \theta(i)}{n} = \frac{\theta(n)(n+1)}{n} = \frac{\theta\left(\frac{n^2+n}{2}\right)}{n} = \theta(n)$$

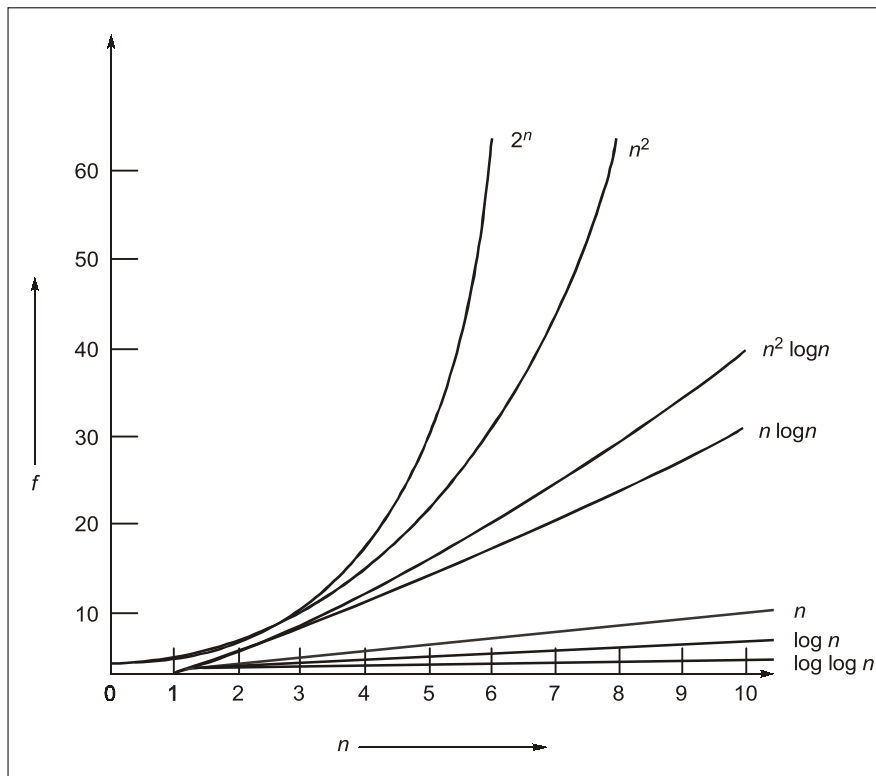
### 1.2.3 Best Case Analysis

In the best case analysis, we calculate lower bound on running time of an algorithm. We must know the case that causes minimum number of operations to be executed. In the linear search problem, the best case occurs when  $x$  is present at the first location. The number of operations in the best case is constant (not dependent on  $n$ ). So time complexity in the best case would be  $\Theta(1)$ .

- Most of the times, we do worst case analysis to analyze algorithms. In the worst analysis, we guarantee an upper bound on the running time of an algorithm which is good information.
- The average case analysis is not easy to do in most of the practical cases and it is rarely done. In the average case analysis, we must know (or predict) the mathematical distribution of all possible inputs.
- The Best Case analysis is bogus. Guaranteeing a lower bound on an algorithm doesn't provide any information as in the worst case, an algorithm may take years to run.

## 1.3 Asymptotic Notations

Let  $f$  be a non negative function. Then we can define the three most common asymptotic bounds as follows.



$$O(1) < O(\log_2 n) < O(n) < O(n^2) < O(n^3) \dots < O(n^k) < O(2^n)$$

### Summary



- In the worst case analysis, we calculate upper bound on running time of an algorithm.
- In average case analysis, we take all possible inputs and calculate computing time for all of the inputs.
- In the best case analysis, we calculate lower bound on running time of an algorithm.
- We say that  $f(n)$  is Big-O of  $g(n)$ , written as  $f(n) = O(g(n))$
- We say that  $f(n)$  is Big-Omega of  $g(n)$ , written as  $f(n) = \Omega(g(n))$
- We say that  $f(n)$  is **Big-Theta** of  $g(n)$ , written as  $f(n) = \Theta(g(n))$
- $O(1) < O(\log_2 n) < O(n) < O(n^2) < O(n^3) \dots < O(n^k) < O(2^n)$
- $\log(n!) < (\log(n))!$
- $\log(\log^* n) < \log^*(\log n)$
- $n < (\log n)^{\log n}$
- $2^n < n! < n^n$



**Student's  
Assignments**

**1**

**Q.1** Which one of the following is true?

1.  $an = o(n^2)$   $a \geq 0$
2.  $an^2 = O(n^2)$   $a > 0$
3.  $an^2 \neq o(n^2)$   $a > 0$

- (a) Only 1 and 2 are correct
- (b) Only 1 is correct
- (c) 1 and 3 are correct only
- (d) All are correct

**Q.2** Which of the following statements is true?

**S1:**  $\frac{1}{2}n^2 = \omega(n)$       **S2:**  $\frac{1}{2}n^2 = \omega(n^2)$

- (a) S1 is correct
- (b) S2 is correct
- (c) S1 and S2 both are correct
- (d) None of the above

**Q.3**  $f(n) = 3n^2 + 4n + 2$ . Which will be the exact value for  $f(n)$

- (a)  $\Theta(n^2)$
- (b)  $o(n^2)$
- (c)  $O(n^2)$
- (d)  $\Omega(n^2)$

**Q.4**  $f(n) = O(g(n))$  if and only if

- (a)  $g(n) = O(f(n))$
- (b)  $g(n) = \omega(f(n))$
- (c)  $g(n) = \Omega(f(n))$
- (d) None of these

**Q.5**  $f(n) = o(g(n))$  if and only if

- (a)  $g(n) = \Omega(f(n))$
- (b)  $g(n) = \omega(f(n))$
- (c)  $g(n) = \Omega(f(n))$  and  $g(n) = \omega(f(n))$
- (d) None of these

**Q.6** Which of the following is not correct?

- (a)  $f(n) = O(f(n))$
- (b)  $c \cdot f(n) = O(f(n))$  for a constant  $C$
- (c)  $(f(n) + g(n)) = o(g(n) + f(n))$
- (d) None of the above

**Q.7**  $f(n) = \Theta(g(n))$  is

- (a)  $0 \leq C_1 g(n) \leq f(n) \leq C_2 g(n) \forall n \geq n_0$  where  $C_1, C_2, n_0$  are + integer
- (b)  $0 \leq C_1 g(n) \leq f(n) \forall n \geq n_0$  where  $C_1, C_2, n_0$  are + integer

- (c)  $0 \leq f(n) \leq C g(n) \forall n \geq n_0$  where  $C_1, n_0$  are + integer
- (d) None of the above

**Q.8**  $f(n) = O(g(n))$  implies

- (a)  $0 \leq C_1 g(n) \leq f(n) \forall n \geq n_0$  where  $C_1, n_0$  are + integer
- (b)  $0 \leq C_1 f(n) \leq C_2 g(n) \forall n \geq n_0$  where  $C_1, C_2, n_0$  are + integer
- (c)  $0 \leq f(n) \leq C g(n) \forall n \geq n_0$  where  $C, n_0$  are + integer
- (d)  $0 \leq C_1 g(n) \leq C_2 f(n) \forall n \geq n_0$  where  $C_1, C_2, n_0$  are + integer

**Q.9**  $f(n) = \Theta(g(n))$  implies

- (a)  $f(n) = O(g(n))$  only
- (b)  $f(n) = \Omega(g(n))$  only
- (c)  $f(n) = O(g(n))$  and  $f(n) = \Omega(g(n))$
- (d) None of the above

**Q.10**  $f(n) = o(g(n))$  implies

- (a)  $0 \leq f(n) \leq Cg(n)$  such that there exists some positive constant  $C$  and  $n_0 > \forall n \geq n_0$
- (b)  $0 \leq f(n) < Cg(n)$  for every +ve constant  $C > 0$  there exists  $n_0 > 0, \forall n \geq n_0$
- (c)  $0 \leq C_1 f(n) \leq C_2 g(n) \forall n \geq n_0$  such that  $C_1, C_2$  and  $n_0$  are +ve constants
- (d)  $0 \leq f(n) < Cg(n)$  for some +ve constant  $C > 0 \exists n_0 > 0, \forall n \geq n_0$

**Q.11** void x (int A[ ], int n)

```
{
    int i, j;
    for (i = 0; i < n; i++)
    {
        j = n - 1;
        while (j > i)
        {
            swap (A[j], A[j - 1]);
            j--;
        }
    }
}
```

What will be time complexity of the above algorithm if swap function takes constant time?

- (a)  $O(n)$
- (b)  $O(n^2)$
- (c)  $O(n \log_2 n)$
- (d)  $O(n^3)$

**Q.12** Which of the following statements are True?

- (a)  $100 n \log n = O(n \log n / 100)$
- (b)  $\sqrt{\log n} = O(\log \log n)$
- (c) If  $0 < x < y$  then  $n^x = O(n^y)$
- (d)  $2^n \neq O(n^c)$  where  $c$  is a constant and  $c > 0$

**Q.13** Which of the following statements ( $k, m$  are constants) are True?

- (a)  $(n + k)^m = O(n^m)$
- (b)  $2^{n+1} = O(2^n)$
- (c)  $2^{2^n} = O(2^n)$
- (d)  $f(n) = O(f(n)^2)$

**Q.14** Which of the following statements are True?

- (a)  $n^2 \cdot 2^{3 \log_2 n} = \Theta(n^5)$
- (b)  $\frac{4^n}{2^n} = \Theta(2^n)$
- (c)  $2^{\log_2 n} = \Theta(n^2)$
- (d) if  $f(n) = O(g(n))$  then  $2^{f(n)} = O(2^{g(n)})$

**Q.15** Consider  $f(n)$ ,  $g(n)$  and  $h(n)$  be function defined as follows:

$$\begin{aligned} f(n) &= \Omega(n^3) \\ g(n) &= O(n^2) \\ h(n) &= \Theta(n^2) \end{aligned}$$

Which of the following represents correct asymptotic solution for  $f(n) + [g(n) \times h(n)]$ ?

- (a)  $\Omega(n^3)$
- (b)  $O(n^4)$
- (c)  $\Theta(n^4)$
- (d)  $O(n^3)$

**Q.16** Consider the following C-function:

```
int Rec (int n)
{
    int i, j, k, p, q = 0;
    for (i = n; i > 1; i/2) {
        p = 0;
        for (j = 1; j < n; j++)
            p = p + 1;
        for (k = 1; k < p; k = k * 3)
            q++;
    } return q
}
```

Then time complexity of Rec in term of  $\Theta$  notation is

- (a)  $\Theta(n)$
- (b)  $\Theta(n^2)$
- (c)  $\Theta(n \log n)$
- (d)  $\Theta(n \log \log n)$

**Q.17** Consider following functions:

$$\begin{aligned} f(n) &= (\log n)^{n-1} \\ g(n) &= 2^n \\ h(n) &= e^n / n \end{aligned}$$

Which of the following is true?

- (a)  $f(n) = O(h(n))$
- (b)  $f(n) \neq O(n(n))$
- (c)  $g(n) = \Omega(f(n))$
- (d) None of these

**Q.18** Consider the following C function:

```
void dosomething (int n)
{
    int m, j, k;
    for (j = 0; j < 200; j++)
    {
        for (k = 0; k < n; k++)
        {
            for (m = 0; m < j; m++)
                printf("%d", i + m);
        }
    }
}
```

What is the time complexity of the above function?

- (a)  $O(n^2)$
- (b)  $O(n \log n)$
- (c)  $O(n)$
- (d)  $O(n^2 \log n)$

**Q.19** Let  $g(n) = \Omega(n)$ ,  $f(n) = O(n)$  and  $h(n) = \Theta(n)$  then what is the time complexity of  $[g(n) \cdot f(n) + h(n)]$ ?

- (a)  $O(n)$
- (b)  $\Theta(n)$
- (c)  $\Omega(n)$
- (d)  $\Theta(n^2)$

**Q.20** Consider the following functions:

$$f_1 = n^4, f_2 = 4^n, f_3 = n^{110/37}, f_4 = \left(\frac{119}{37}\right)^n$$

Which of the following is correct order of increasing growth rate?

- (a)  $f_1, f_3, f_2, f_4$
- (b)  $f_3, f_1, f_4, f_2$
- (c)  $f_3, f_1, f_2, f_4$
- (d)  $f_1, f_3, f_4, f_2$

**Q.21** Consider the following program segments:

```
main ( )
{
    int i = 0, m = 0;
    for (i = 1; i < n; i++) {
        for (j = 1; j < i * i; j++) {
            if ((j % i) == 0)
                for (k = 1; k < j; k++) {
```

```

        m = m + 1; }
    }
}

```

What is the time complexity of above program?

- (a)  $O(n^2 \log n)$  (b)  $O(n^3)$   
(c)  $O(n^4)$  (d)  $O(n \log n)$

**Answer Key:**

1. (d) 2. (a) 3. (a) 4. (c) 5. (c)  
6. (c) 7. (a) 8. (c) 9. (c) 10. (b)  
11. (b) 12. (a, c, d) 13. (a, b) 14. (a, b) 15. (a)  
16. (c) 17. (b) 18. (c) 19. (c) 20. (b)  
21. (c)



**Student's  
Assignments**

**Explanations**

1. (d)

1.  $an = o(n^2)$ ,  $a \geq 0$

$0 \leq an < c \cdot n^2$  for all value of  $c$  this condition is true. Hence 1 is correct.

2.  $an^2 = O(n^2)$

$0 \leq an^2 \leq c \cdot n^2$  when  $c \geq a$  this condition is true. Hence, 2 is correct.

3.  $an^2 \neq o(n^2)$

for this statement to be false  $0 \leq an^2 < c \cdot n^2$  must be true

Since above condition needs to be true for all  $c$  thus if we could prove for some  $c$  that violates above condition. Whole option becomes true.  
 $0 \leq an^2 < c \cdot n^2$  is not true when  $c < a$  thus  $an^2 \neq o(n^2)$  is true.

2. (a)

**S1:**  $c \cdot \frac{1}{2}n^2 > n$  for every  $c > 0$  which can be seen easily.

**S2:** False since  $c \cdot \frac{1}{2}n^2 > n^2$  is false when  $c = 2$  hence for every  $c$  it is not true thus, false.

3. (a)

$$f(n) = 3n^2 + 4n + 2$$

maximum degree = 2

$$\Rightarrow \theta(n^2)$$

**Note:** option (c) and (d) are also possible but they are not exact.

4. (c)

$$f(n) = O(g(n))$$

Take some arbitrary values in order to solve these type of questions.

As given,  $f(n) = O(g(n))$

$$\Rightarrow f(n) \leq g(n)$$

Let  $f(n) = n$  and  $g(n) = n^2$

(a)  $g(n) = O(f(n)) \rightarrow n^2 = O(n)$  or  $n^2 \leq n$  not true

(b)  $g(n) = \omega(f(n)) \rightarrow g(n) > f(n) \rightarrow n^2 > n$

True (only for assumed value).

But it become false when both  $f$  and  $g$  are  $n$ .

(c)  $g(n) = \Omega(f(n)) \rightarrow g(n) \geq f(n)$

Same thing written as given in question.

5. (c)

$$f(n) = o(g(n))$$

$$\Rightarrow f(n) \text{ is strictly lesser than } g(n)$$

Thus,  $f(n) < g(n)$

For (c),  $g(n) = \Omega(f(n))$

$$\Rightarrow g(n) \geq f(n)$$

Satisfied by given equation.

$$\Rightarrow g(n) = \omega(f(n))$$

$$\Rightarrow g(n) > f(n) \text{ this is also satisfied}$$

6. (c)

(a)  $f(n) = O(f(n))$  by reflexive property.

(b)  $c * f(n) = O(f(n))$  since constant does not matter in case of asymptotic analysis.

(c)  $(f(n) + g(n)) = o(g(n) + f(n))$

let  $f(n) = n^2$

$$g(n) = n$$

$$(n^2 + n) = o(n^2 + n)$$

$$n^2 = o(n^2) \text{ not true since tightest bound not allowed.}$$

9. (c)

Go Through the Topic  $\theta$  notation.

$$F(n) = \theta(g(n)) \text{ if and only if } f(n) = O(g(n)) \text{ and}$$

$$F(n) = \Omega(g(n)) \text{ simultaneously.}$$

10. (b)

As Small-oh definition says

$$f(n) = o(g(n))$$

when  $0 \leq f(n) < c \cdot g(n)$  for every  $c > 0$  and for all  $n > n_0$

This is the major point.

Go Through the difference between Big-oh and Small-oh notation for this concept.

Option (d) is false since for some +ve constant is written there.

11. (b)

When  $i = 0$  $j$  loop i.e. while loop runs for  $(n - 1)$  timewhen  $i = 1$  $j$  loop i.e. while loop runs for  $(n - 2)$  time

⋮

when  $i = n - 1$  $j$  loop will not runHence,  $TC = 1 + 2 + \dots (n - 1)$ 

$$= \theta \left\{ \frac{(n-1)(n)}{2} \right\}$$

i.e.  $\theta(n^2)$ 

12. (a, c, d)

(a)  $100 n \log n = O(n \log n / 100)$ 

constant does not matter. So it becomes

$$n \log n = O(n \log n)$$

Reflexive property. (True)

(b)  $\sqrt{\log n} = O(\log \log n)$ 

$$\Rightarrow \sqrt{\log n} \leq \log \log n$$

Taking log on both sides

$$\frac{1}{2} \underbrace{\log \log n}_x \leq \log(\log \log n)$$

$x \leq \log x$  but this is not valid since  $x > \log x$  thus, false statement.

(c) If  $0 < x < y$  then  $n^x = O(n^y)$ 

Exponential grows faster than polynomial.

Hence, true.

(d)  $2^n \neq O(n^c)$  since  $n^c = O(2^n)$  same reason as of  $c$  option.

13. (a, b)

(a)  $(n + k)^m = O(n^m)$ Take  $m = 2$  and  $k = 1$ 

$$(n + 1)^2 = O(n^2)$$

$$n^2 + 1 + 2n = O(n^2) \quad (\text{True})$$

(b)  $2^n \cdot 2 = O(2^n)$  true since constant does not matter.

(c)  $2^{2^n} = O(2^n)$ Let,  $2^n = x$ 

$$2^x = O(x) \quad (\text{False})$$

Exponential grows faster.

(d)  $f(n) = O(f(n)^2)$ Since we know  $x > x^2$  when  $x < 1$  thus,

$$\text{Taking } f(n) = \frac{1}{n}$$

$$\Rightarrow \frac{1}{n} = O\left(\frac{1}{n^2}\right)$$

$$\text{False since } \frac{1}{n} > \frac{1}{n^2}$$

14. (a, b)

(a)  $n^2 \cdot 2^{3 \log_2 n} = \theta(n^5)$ 

$$\text{Note: } x^{\log_x n} \Leftrightarrow n$$

$$\log_2 n^x \Leftrightarrow x \log_2 n$$

$$\begin{aligned} \text{Thus, } n^2 \cdot 2^{\log_2 n^3} \\ n^2 \cdot n^3 \\ n^5 = \text{RHS} \quad (\text{True}) \end{aligned}$$

$$(b) \frac{4^n}{2^n} = \left(\frac{4}{2}\right)^n = 2^n = \theta(2^n) \quad (\text{True})$$

(c)  $2^{\log_2 n} = \theta(n^2)$   
 $n = \theta(n^2)$  not true since,  $n \neq \Omega(n^2)$

(d) If  $f(n) = O(g(n))$  then  $2^{f(n)} = O(2^{g(n)})$ Let  $f(n) = 2n$ 

$$g(n) = n$$

$$\Rightarrow 2n = O(n) \quad (\text{True})$$

But  $2^{2n} = O(2^n)$  false since for this to be true  $c$  must be  $2^n$  but  $c$  cannot be any function since it is constant.

15. (a)

•

$$f(n) = \Omega(n^3)$$

$$f(n) \geq C_1(n^3)$$



Q.5 Arrange the following in increasing order:

$$f_1 = \left(\frac{3}{2}\right)^n, f_2 = n^3, f_3 = \log_2 n,$$

$$f_4 = \log(n!), f_5 = 2^{2^n}, f_6 = n^{1/\log n}$$

- (a)  $f_1 f_2 f_3 f_4 f_5 f_6$  (b)  $f_6 f_3 f_4 f_2 f_5 f_1$   
(c)  $f_6 f_4 f_3 f_2 f_1 f_5$  (d)  $f_6 f_3 f_4 f_2 f_1 f_5$

Q.6 Indicate for each pair of expression (A, B) in the table below whether A is O, o,  $\Omega$ ,  $\omega$  or  $\theta$  of B. Fill the table with Yes or No.

A	B	O	o	$\Omega$	$\omega$	$\theta$
$\log^k n$	$n^E$					
$n^k$	$c^n$					
$\sqrt{n}$	$n^{\sin}$					
$n^{\log c}$	$c^{\log n}$					
$\log(n!)$	$\log(n^n)$					

Q.7 Consider the following functions:

$$f(n) = 2^{\log_2 n}$$

$$g(n) = n^{\log_2 n}$$

$$h(n) = n^{1/\log_2 n}$$

Which of the following statements about the asymptotic behaviour of  $f(n)$ ,  $g(n)$  and  $h(n)$  is true?

- (a)  $f(n) = \Omega(g(n))$  and  $g(n) = O(h(n))$   
(b)  $g(n) = \Omega(h(n))$  and  $f(n) = O(f(n))$   
(c)  $f(n) = O(g(n))$  and  $g(n) = \Omega(h(n))$   
(d)  $g(n) = O(h(n))$  and  $h(n) = O(g(n))$

Q.8 Let  $f(n) = \Omega(n)$ ,  $g(n) = O(n)$  and  $h(n) = \theta(n)$ . Then  $[f(n) \cdot g(n)] + h(n)$  is \_\_\_\_\_.

- (a)  $\Omega(n)$  (b)  $O(n)$   
(c)  $\theta(n)$  (d) None of these

Q.9 What is time complexity of fun ( )?

```
int fun (int n)
{
    int count = 0;
    for (int i = n; i > 0; i/=2)
        for (int j = 0; j < i; j++)
            count += 1;
    return count;
}
```

(a)  $O(n^2)$  (b)  $O(n)$   
(c)  $O(n \log n)$  (d)  $O(n(\log n)^2)$

Q.10 What is the time complexity of the following program?

```
void fun(int n) {
    int i, j, counter = 0;
    for (i = 1; i ≤ n; i++) {
        for (j = 1; j × j ≤ n; j++)
            count ++;
    }
}
```

- (a)  $O(n \log n)$  (b)  $O(n^2)$   
(c)  $O(n \log \log n)$  (d)  $O(n^{3/2})$

Q.11 Consider the following function:

<pre>void f(int n) {     int i, j, k, m;     for (i = 0; i &lt; 100; i++)     {         for (j = 0; j &lt; n; j++)         {             for (k = 0; k &lt; j; k++)                 g(j, k);             for (m = 0; m &lt; i; m++)                 g(k, m);         }     } }</pre>	<pre>void g(int a, int b) {     int temp;     temp = a;     a = b;     b = temp; }</pre>
--	--

What is the worst case running time of the function  $f$  for any positive value of  $n$ ?

- (a)  $O(1)$  (b)  $O(n)$   
(c)  $O(n^2)$  (d)  $O(n^3)$

Q.12 Which of the following is correct?

- (a)  $(2^{\log_2 \log_2 n})^2 < 2^{(\log_2 \log_2 n)^2} < 2^{\sqrt{\log_2 n}} < (\log_2 n)^{\log_2 n}$   
(b)  $2^{(\log_2 \log_2 n)^2} < (2^{\log_2 \log_2 n})^2 < 2^{\sqrt{\log_2 n}} < (\log_2 n)^{\log_2 n}$   
(c)  $(2^{\log_2 \log_2 n})^2 < 2^{\sqrt{\log_2 n}} < 2^{(\log_2 \log_2 n)^2} < (\log_2 n)^{\log_2 n}$   
(d) None of the above

Q.13 Consider the following statement:

$$S_1: f(n) = O(f(n)^2)$$

$$S_2: \text{If } f(n) = \Omega(g(n)) \text{ then } f(n) = O(g(n))$$

$$S_3: \text{If } f(n) \neq O(g(n)) \text{ then } g(n) \neq O(f(n))$$

Which of the following statements are always true?

- (a) Only  $S_1$  and  $S_3$  (b) Only  $S_1$  and  $S_2$   
(c) Only  $S_2$  and  $S_3$  (d) None of these



**Answer Key:**

1. (c)    2. (a)    3. (b)    4. (a)    5. (d)  
6. Sol.    7. (c)    8. (a)    9. (b)    10. (d)  
11. (c)    12. (a)    13. (d)



**Student's  
Assignments**

**Explanations**

1. (c)

$n! = n \times n-1 \times n-2 \times \dots \times 2 \times 1$   
If we keep smaller value at every place  
i.e.  $2 \times 2 \times 2 \times \dots \times 2 \times 2$   
Intuitively  $n! > 2^n$   
Thus,  $n! = \omega(2^n)$   
If we keep largest value at every place i.e.  
 $n \times n \times n \times n \dots \times n \times n$   
Then,  $n^n > n!$   
Thus,  $n^n > \omega(n^n)$   
Both I and II are true.

2. (a)

When,  $n = 2^x$   
 $\log^* n = \log^* 2^x = (1 + \log^* x)$   
 $\log n = \log 2^x = x$   
Let  $\log(\log^* n) < \log^*(\log n)$   
If this is true this means

$$\lim_{n \rightarrow \infty} \frac{\log(\log^* n)}{\log^*(\log n)} = 0$$

$$\lim_{x \rightarrow \infty} \frac{\log(1 + \log^* x)}{\log^* x}$$

Let  $\log^* x = y$

$$\lim_{y \rightarrow \infty} \frac{\log(1+y)}{y} \left( \frac{\infty}{\infty} \text{ form} \right)$$

$$\text{L hospital } \lim_{y \rightarrow \infty} \frac{1}{1+y} = 0$$

Hence,  $f(n) = O(g(n))$

3. (d)

(a) Let  $f(n) = n$ ,  $g(n) = n^2$

$$\text{Thus } \underbrace{f(n) = O(g(n))}_{\text{True}} \rightarrow \underbrace{g(n) = O(f(n))}_{\text{False}}$$

(b) Given option is the well known method for checking the inequality. We can take log on both sides and nothing will change.

(c)  $f(n) + O(f(n)) = \theta(f(n))$

Strictly greater than  $f(n)$

Thus LHS =  $\theta$  (RHS) can never be true.

Hence given statement is not true.

(d)  $f(n) = 2^n$

Then,  $2^n = \theta(2^{n/2})$  (Not true)

Hence, number of correct statements = 1.

4. (a)

$$f(n) = n^{\frac{2 + \cos n}{[-1, 1]}}$$

Thus,  $f(n) = [n, n^3]$

$$g(n) = n^{\frac{\sin n}{[-1, 1]}} = [n^{-1}, n]$$

$$\Rightarrow g(n) = O(f(n))$$

5. (d)

$$f_5 = 2^{2^n} f_6 = n^{1/\log n}$$

Let  $2^{2^n} < n^{1/\log n}$

Taking log on both side,

$$2^n < \frac{1}{\log n} \cdot \log n \quad (\text{Not true thus, } f_5 > f_6)$$

$f_1$  is exponential, taking  $f_1$  and  $f_6$

$$\text{Let } \left(\frac{3}{2}\right)^n < n^{1/\log n}$$

$$n \log_2 3 < \frac{1}{\log n} \cdot \log n \quad (\text{Not true thus } f_1 > f_6)$$

$f_4 = (n \log n)$  Stirling approximation

$$f_6 < f_3 < f_4 < f_2 < f_1 < f_5$$

6. Sol.

A	B	O	o	$\Omega$	$\omega$	$\theta$
$\log^K n$	$n^E$	Yes	Yes	No	No	No
$n^k$	$c^n$	Yes	Yes	No	No	No
$\sqrt{n}$	$n^{\sin}$	Yes	No	Yes	No	Yes
$n^{\log c}$	$c^{\log n}$	No	No	Yes	Yes	No
$\log(n!)$	$\log(n^n)$	Yes	Yes	No	No	No

Polynomial and exponential

$\log(n!) < \log(n^n)$

Polynomial grows faster than logarithm

Let  $n^{1/2} \leq n^{\sin n}$

$$\frac{1}{2} \log n \leq \frac{\sin n}{\text{constant}} (\log n)$$

7. (c)

$$\begin{aligned} f(n) &= 2^{\log_2 n} = n^{\log_2 2} = n \\ g(n) &= n^{\log n} \\ h(n) &= n^{1/\log n} \end{aligned}$$

$$= {}^{\log n}\sqrt{n} \left[ n > {}^{\log n}\sqrt{n} \text{ for all large value of } n \right]$$

[It is less than  $n$  since max power of  $n$  is always less than 1 for large value of  $n$ ]

So,  $g(n) \geq f(n) \geq h(n)$

So,  $f(n) = O(g(n))$  and  $g(n) = \Omega(h(n))$

8. (a)

$$\begin{aligned} f(n) &= \Omega(n) \Rightarrow f(n) \geq c \cdot n \\ g(n) &= O(n) \Rightarrow g(n) \leq c \cdot n \\ h(n) &= \theta(n) \Rightarrow c_1 \cdot n \leq h(n) \leq c_2 \cdot n \\ f(n) \cdot g(n) &= c \cdot n \quad [\because f(n) \geq c \cdot n \text{ \& } g(n) \leq c \cdot n] \\ \frac{f(n) \cdot g(n)}{\geq c \cdot n} + \frac{h(n)}{\theta(n)} &= \Omega(n) \end{aligned}$$

So, option (a) is correct.

9. (b)

For  $n$  time, inner loop will execute for  $n$  times.

For  $\frac{n}{2}$  time, inner loop will execute for  $\frac{n}{2}$  times.

For  $\frac{n}{4}$  time, inner loop will execute for  $\frac{n}{4}$  times.

and do on .....

So time complexity:

$$T(n) = O\left(n + \frac{n}{2} + \frac{n}{4} + \dots + 1\right) = O(n)$$

10. (d)

for ( $i = 1; i \leq n; i++$ )  $\Rightarrow O(n)$

for ( $j = 1; j \times j \leq n; j++$ )  $\Rightarrow O(\sqrt{n})$

count ++;

Total time complexity =  $O(n \times n^{1/2}) = O(n^{3/2})$

11. (c)

$$\begin{aligned} f(n) &= \sum_{i=0}^{99} \sum_{j=0}^{n-1} \left( \sum_{k=0}^{j-1} 1 + \sum_{m=0}^{i-1} 1 \right) \\ &= \sum_{i=0}^{99} \sum_{j=0}^{n-1} \sum_{k=0}^{j-1} 1 + \sum_{i=0}^{99} \sum_{j=0}^{n-1} \sum_{m=0}^{i-1} 1 \\ &= O(n^2) + O(n) = O(n^2) \end{aligned}$$

12. (a)

$2^{\sqrt{\log_2 n}}$  growth rate is fast.

Since, for bigger values in above table sequence:

$$(2^{\log_2 \log_2 n})^2 < 2^{(\log_2 \log_2 n)^2} < 2^{\sqrt{\log_2 n}} < (\log_2 n)^{\log_2 n}$$

13. (d)

$$S_1: f(n) = O(f(n)^2)$$

$$\text{If } f(n) = 1/n$$

$$\text{then } f(n) \neq O(f(n)^2)$$

$$\text{since } 1/n \geq C_1(1/n)^2$$

So, false, i.e., not always **true**.

$$S_2: f(n) = \Omega(g(n)) \text{ and } f(n) = O(g(n))$$

$$\text{Since } f(n) = n \text{ and } g(n) = n^2$$

$$\text{then } f(n) = O(g(n)) \text{ but not } f(n) = \Omega(g(n))$$

So, false i.e., not always **true**.

$$S_3: f(n) \neq O(g(n)) \text{ and } g(n) \neq O(f(n))$$

$$\text{If } f(n) = n^2 \text{ and } g(n) = n$$

$$\text{then } f(n) \neq O(g(n)) \text{ and } g(n) = O(f(n))$$

i.e., not always true.

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